

Recall the standard primal-dual pair of CLPs:

$$(P) \quad v_p^* = \inf \langle c, x \rangle$$

$$\text{s.t. } \langle a_i, x \rangle = b_i, \quad i=1, \dots, m,$$

$$x \in K \subseteq E \quad (x \succ_K 0)$$

$$(D) \quad v_d^* = \sup \sum_{i=1}^m b_i y_i \quad (= b^T y)$$

$$\text{s.t. } c - \sum_{i=1}^m y_i a_i \in K^* \quad (s \succ_{K^*} 0)$$

Here, E is an Euclidean space, K is a closed pointed cone with non-empty interior in E , $c, a_i \in E$, $b_i \in \mathbb{R}$, $\langle \cdot, \cdot \rangle$ is an inner product on E ,

$$K^* = \{ w \in E : \langle w, x \rangle \geq 0 \quad \forall x \in K \}$$

is the dual cone of K ,

Example (Eigenvalue optimization)

Let $A_1, \dots, A_k \in S^n$ be given symmetric matrices.

Goal - Find $y \in \mathbb{R}^k$ to minimize $\lambda_{\max} \left(\sum_{i=1}^k y_i A_i \right)$.

Proposition: Let $A \in S^n$ be given. Then,

$$\lambda_{\max}(A) \leq t \iff tI \succ_{S_+^n} A.$$

(i.e., $tI - A \in S_+^n$)

With the above proposition, we have

$$\min_y \lambda_{\max} \left(\sum_{i=1}^k y_i A_i \right) \iff \min_{y, t} t \iff \min_{y, t} t$$

$$\text{s.t. } t \geq \lambda_{\max} \left(\sum_{i=1}^k y_i A_i \right) \quad \text{s.t. } tI - \sum_{i=1}^k y_i A_i \in S_+^n$$

↑
SDP

Proof of Proposition:

(\Leftarrow) Suppose that $tI - A \in S_+^n$. By definition, for all $u \in \mathbb{R}^n, \|u\|_2 = 1$,

$u^T(tI - A)u \geq 0$. This implies that

$$t \cancel{\|u\|_2^2}^1 \geq u^T A u \Rightarrow t \geq u^T A u \Rightarrow \max_{\|u\|_2=1} u^T A u \leq t$$

$\Rightarrow \lambda_{\max}(A) \leq t$ (by Courant-Fischer theorem)

(\Rightarrow) Reverse the above argument

Theorem (Weak Duality for LP)

Let \bar{x} be feasible for (P), \bar{y} be feasible for (D). Then,

$$b^T \bar{y} \leq \langle c, \bar{x} \rangle.$$

Proof: We compute

$$\begin{aligned} \langle c, \bar{x} \rangle - b^T \bar{y} &= \langle c, \bar{x} \rangle - \sum_{i=1}^m \bar{y}_i \underbrace{\langle a_i, \bar{x} \rangle}_{= b_i} \\ &\quad (\because \bar{x} \text{ feasible for (P)}) \\ &= \langle c, \bar{x} \rangle - \left\langle \sum_{i=1}^m \bar{y}_i a_i, \bar{x} \right\rangle \quad (\text{by linearity of inner product}) \\ &= \left\langle c - \sum_{i=1}^m \bar{y}_i a_i, \bar{x} \right\rangle \\ &\geq 0 \quad (\because \bar{x} \in K, c - \sum_{i=1}^m \bar{y}_i a_i \in K^*) \end{aligned}$$

Theorem (Strong Duality for LP)

Suppose that (P) is bounded below and satisfies Slater's condition; i.e., (D) is bounded above

there exists a feasible \bar{x} for (P) s.t. $\bar{x} \in \text{int}(K)$
 \bar{y} for (D) s.t. $c - \sum_{i=1}^m \bar{y}_i a_i \in \text{int}(K^*)$.

Then, $V_p^* = V_d^*$ (Zero duality gap) and there exists an optimal

dual solution y^* s.t. $b^T y^* = V_p^* = V_d^*$.
 primal solution x^* $\langle c, x^* \rangle = V_p^* = V_d^*$.

Note: In LP strong duality, we do not need to assume Slater condition
 Also, if the duality gap is zero, then both primal and dual have optimal solutions.

Example: Consider the SOCP

$$0 = V_d^* = \sup -y_1$$

$$\hookrightarrow \sup -e_1^T y \quad (e_1 = (1, 0))$$

1

$$0 = V_d^* = \sup -y_1$$

↓

$$y_1 = \frac{1}{4}y_2$$

Take $y_2 \nearrow \infty$

$$\Rightarrow y_1 \downarrow 0$$

However, there is no optimal solution.

$$\text{s.t. } (y_1 + y_2, 1, y_1 - y_2) \in \mathcal{Q}^3$$

↕

$$y_1 + y_2 \geq \sqrt{1^2 + (y_1 - y_2)^2}$$

↕

$$4y_1 y_2 \geq 1, y_1 + y_2 > 0$$

$$(\Rightarrow y_1, y_2 > 0)$$

Q: $\exists \bar{y}_1, \bar{y}_2 > 0$ s.t.

$$\bar{y}_1 + \bar{y}_2 > \sqrt{1 + (\bar{y}_1 - \bar{y}_2)^2}?$$

A: Yes, e.g., $\bar{y}_1 = \bar{y}_2 = 1$ Thus, (1,1)

satisfies Slater's condition.

$(e_1 = (1, 0))$

$$\Leftrightarrow \sup \underbrace{-e_1^T y}_b$$

$$\text{s.t. } \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} - y_1 \begin{bmatrix} -1 \\ 0 \\ -1 \end{bmatrix} - y_2 \begin{bmatrix} -1 \\ 0 \\ 1 \end{bmatrix} \in \mathcal{Q}^3$$

$\underbrace{\quad}_c \quad \underbrace{\quad}_{a_1} \quad \underbrace{\quad}_{a_2}$

$$\sup b^T y$$

$$\text{s.t. } c - \sum_{i=1}^2 y_i a_i \in \mathcal{Q}^3$$

↕ Take the dual

$$0 = V_p^* = \inf \langle c, x \rangle = x_2$$

$$\text{s.t. } \langle a_1, x \rangle = -x_1 - x_3 = b_1 = -1,$$

$$\langle a_2, x \rangle = -x_1 + x_3 = b_2 = 0,$$

$$x = (x_1, x_2, x_3) \in \mathcal{Q}^3$$

This has only one feasible (and hence optimal) solution: $(\frac{1}{2}, 0, \frac{1}{2})$.