

Recall that the SDR technique applies to both real and complex QCPs:

$$\begin{array}{ll}
 (Q_R) & \inf \quad x^T C x \\
 & \text{s.t.} \quad x^T Q_i x \geq b_i, \quad i=1, \dots, m. \\
 & \quad \quad x \in \mathbb{R}^n
 \end{array}
 \quad \text{Data:}
 \begin{array}{l}
 C, Q_1, \dots, Q_m \in \mathcal{S}^n \\
 b_1, \dots, b_m \in \mathbb{R}
 \end{array}$$

$$\begin{array}{ll}
 (SDR_R) & \inf \quad \langle C, X \rangle \\
 & \text{s.t.} \quad \langle Q_i, X \rangle \geq b_i, \\
 & \quad \quad X \in \mathcal{S}_+^n
 \end{array}$$

$$\begin{array}{ll}
 (Q_C) & \inf \quad z^H C z \\
 & \text{s.t.} \quad z^H Q_i z \geq b_i, \quad i=1, \dots, m. \\
 & \quad \quad z \in \mathbb{C}^n
 \end{array}
 \quad \text{Data:}
 \begin{array}{l}
 C, Q_1, \dots, Q_m \in \mathcal{H}^n \\
 b_1, \dots, b_m \in \mathbb{R}
 \end{array}$$

$$\begin{array}{ll}
 (SDR_C) & \inf \quad \langle C, Z \rangle \\
 & \text{s.t.} \quad \langle Q_i, Z \rangle \geq b_i, \quad \rightarrow z^* \text{ opt. soln.} \\
 & \quad \quad z \in \mathcal{H}_+^n, \quad \text{rank}(Z) \leq 1
 \end{array}$$

Observe: If the optimal solution to (SDR_C) z^* has $\text{rank} \leq 1$, then it is optimal for (Q_C) .

Q: Are there conditions on (SDR_C) s.t. it has an optimal solution of $\text{rank} \leq 1$? (We say that the relaxation is tight)

Theorem: Consider the SDP

$$\begin{array}{ll}
 (P) & \inf \quad \langle C, Z \rangle \\
 & \text{s.t.} \quad \langle Q_i, Z \rangle \geq b_i, \quad i=1, \dots, m', \\
 & \quad \quad \langle Q_i, Z \rangle = b_i, \quad i=m'+1, \dots, m, \\
 & \quad \quad Z \in \mathcal{H}_+^n.
 \end{array}
 \quad \left. \vphantom{\begin{array}{l} \inf \\ \text{s.t.} \end{array}} \right\} \text{closed feasible set}$$

Suppose that (P) has an optimal solution. Then, there exists an

optimal solution z^* satisfying $\text{rank}(z^*) \leq \lfloor \sqrt{m} \rfloor$. Moreover, z^* can be found efficiently,

Remark: (P) has optimal solution if

- (i) its dual is bounded above and satisfies Slater's condition; or (strong duality of CLP)
- (ii) the feasible set of (P) is compact (closed and bounded) (Weierstrass theorem)

Also, if $m \leq 3$, then (P) has a rank-1 optimal solution.

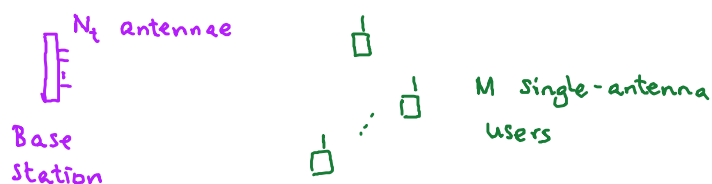
The following is an extension of the above theorem:

Theorem: Consider

$$\begin{aligned}
 \text{(B)} \quad & \inf \sum_{k=1}^K \langle C_k, z_k \rangle \\
 & \text{s.t.} \quad \sum_{k=1}^K \langle Q_{ik}, z_k \rangle \geq b_i, \quad i=1, \dots, m', \\
 & \quad \quad \sum_{k=1}^K \langle Q_{ik}, z_k \rangle = b_i, \quad i=m'+1, \dots, m, \\
 & \quad \quad \underline{z_k} \in \mathbb{H}_+^n, \quad k=1, \dots, K.
 \end{aligned}$$

(Note: (P) is a special case of (B) with $K=1$). Suppose that (B) has an optimal solution. Then, there exists an optimal solution (z_1^*, \dots, z_K^*) satisfying $\sum_{k=1}^K \text{rank}^2(z_k^*) \leq m$.

Example: Unicast transmit beamforming



• Signal transmitted by base station:

$$x(t) = \sum_{i=1}^M s_i(t) w_i,$$

$s_i(t) \in \mathbb{C}$: unit-power symbol ; $w_i \in \mathbb{C}^{N_t}$: beamforming vector for user i :

- signal received by i^{th} user :

$$y_i(t) = h_i^H x(t) + n_i(t)$$

$h_i \in \mathbb{C}^{N_t}$: channel vector of user i
 $n_i(t) \sim \mathcal{CN}(0, \sigma_i^2)$: additive noise

$$= \underbrace{h_i^H w_i s_i(t)}_{\text{signal}} + \underbrace{\sum_{j \neq i} h_i^H w_j s_j(t)}_{\text{interference}} + \underbrace{n_i(t)}_{\text{noise}}$$

- Goal: To minimize transmit power while guaranteeing a certain level of SINR for each user.

Assuming we have the second-order statistics of the channels

$R_i = \mathbb{E}[h_i h_i^H]$, we can formulate the problem as

$$\inf \sum_{i=1}^M \|w_i\|_2^2 = \sum_{i=1}^M w_i^H w_i = \sum_{i=1}^M \langle I, w_i w_i^H \rangle$$

$$\text{s.t. } \text{SINR}_i = \frac{w_i^H R_i w_i}{\sum_{j \neq i} w_j^H R_j w_j + \sigma_i^2} \geq \gamma_i ; \quad i=1, \dots, M.$$

$$w_i \in \mathbb{C}^{N_t}$$

$$\underbrace{w_i^H R_i w_i}_{\text{}} - \gamma_i \underbrace{\sum_{j \neq i} w_j^H R_j w_j}_{\text{}} \geq \gamma_i \sigma_i^2$$

" $\langle R_i, w_i w_i^H \rangle$

The SDR of the above problem is

$$\inf \sum_{i=1}^M \langle I, W_i \rangle$$

$$(U) \quad \text{s.t.} \quad \langle R_i, W_i \rangle - \gamma_i \sum_{j \neq i} \langle R_j, W_j \rangle \geq \gamma_i \sigma_i^2 ; \quad i=1, \dots, M,$$

$$W_1, \dots, W_M \in \mathbb{H}_+^{N_t}$$

Claim: (U) has an optimal solution. (Exercise: Check the dual)

By the second theorem, there exists an optimal solution

(W_1^*, \dots, W_M^*) to (U) s.t. $\sum_{i=1}^M \text{rank}^2(W_i^*) \leq M$. This implies

that $\text{rank}(W_i^*) \leq 1$ if $W_i^* \neq 0 \quad \forall i$. Hence, the SDR is tight.

(verify)